

CS341 #17. Producer Consumer, Semaphores, Condition Variables. Barriers & Reader Writer Problem

1. Producer Consumer & Counting Semaphores (review)

Assume buffer is an array of length 16. Better names for s1?s2?

<pre> 01 void add(value) { 02   sem_wait(&amp;sem_s1) 03   buffer[ (in++) &amp; 15 ] = value; 04   sem_post(&amp;sem_s2); 05 }                 </pre>	<pre> 06 remove() { 07   sem_wait(&amp;sem_s2); 08   result = buffer[ (out++) &amp; 15 ]; 09   sem_post(&amp;sem_s1); 10   return result; 11 }                 </pre>
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Q. What are 'sem\_s1' and sem\_s2? When do they block?

*When internal counter gets to zero*

Q. What should be their initial values?

*16 and 0*

Q. What if sem\_s1 was only initialized to 7? Would the producer consumer still work? to 32?

*Then there are only 7 space available. (Like a train in circle)*

*Overwrite happened !!!*

Q. What is missing from the above code? When would it matter?

*mutex-lock. Without mutex lock, only work when 1 consumer and 1 producer*

Q. Could you implement a producer consumer queue using condition variables instead?

2. Fix the following multithread code to be thread safe, and use condition variables to avoid busy waiting

```

01 #define N (8)
02 pthread_cond_t cvs[N];
03 pthread_mutex_t locks[N];
04 int data[N];
05 int quit;

06 void init() {
07   for(int i =0; i < N;i++) {
08
09     pthread_cond_init(cvs + i, NULL);
10     pthread_mutex_init(locks + i, NULL);
11   }
12 }

13 // Wait until data[i] > 1, then subtract 2 and increment data[i+1]
14 void runner(void*arg) { // For N-1 threads. Each thread gets a value 0 to N-2
15   int i = (int) arg;
16   while(!quit) {
17     while(data[i] < 2) {
18       wait(cvs+i, locks+i) sleep for a bit
19     }
20     data[i] -= 2;
21     data[i+1] ++;
22     } also, broadcast!!!
23 }

24
25 void modify(int index, int amount) {
26   lock(locks[index])
27   data[index] += amount;
28   unlock(locks[index])
29   P_cond - broadcast [cvs+ index]
30   unlock (locks[index])
31 }
                
```

*Binary*

*p\_cond - wait(cvs+i, locks+i)*

*what if p still 2 cookies left!*

3.Counting Semaphore Quick Review I. choose {will always / may / will never} :

sem\_post never block sem\_wait may block.

### 3. Counting Semaphore Quick Review II

10 threads call `sem_wait`. 3 threads immediately continue, the other 7 are blocked. Then `sem_post` is called twice (2). How many additional threads will continue? 2

### 4. Three classic / well known synchronization problems:



Producer Consumer

Reader-Writer Problem

multiple readers  
one exclusive writer

*m-lock won't work!*

*too much wait!*

### 5. pthread barriers

```
pthread_barrier_init( &barrier, 5 );  
pthread_barrier_destroy(&barrier)
```

```
pthread_barrier_wait( &barrier)
```

Return values? *one thread is specially labeled*

0

PTHREAD\_BARRIER\_SERIAL\_THREAD

### 6. Use a CV to implement a single-use barrier until all 8 threads have reached the barrier.

```
int remain = 8
```

```
barrier()
```

*lock*

```
remain -- ;
```

```
if (remain == 0) p_and_broadcast
```

```
else while (remain > 0) { p_cond_wait( &cv, &m ) }
```

*unlock*

*key: another value!*

### 7. Post-lecture challenge:

- Can you make a barrier using only counting semaphores?
- Can you make a barrier using only mutex locks?