> What is a process's umask and how is it used?
What is the default value?
When is it used?
Take away permission from the file created in the future

> Case study: Use mount to explore an iso image
ISO 9660
Example : Use curl –O to download a file
What is an .iso file? → only good for read-only file
CD-ROM

> Starting a virtual machine using an iso file image
What is qemu?

qemu-system-x86_64 -k en-us -cdrom dsl-4.4.10.iso

```
int main() {
  int fd = open("file", O_RD);
  fstat(fd, &s);
  char *ptr = mmap(NULL, s.st_size,
    PROT_READ,
    MAP_FILE | MAP_SHARED, fd, 0);

  for(int i=0; i< s.st_size;i++)
     if(ptr[i] >31) printf("%x %c\n",i, ptr[i]);

}
```

> Welcome to the mmap diner. What would you like?

```
void *
mmap(
  void *addr,
  size_t len,
  int prot,
  int flags,
  int fd,
  off_t off);
```
returns (void*)-1 if failed

Ask yourself -
1. What kind of memory protection would you like?

2. Will the contents of your RAM (random access memory) be backed by a file or will be it anonymous?

3. What happens if you change your RAM contents? Will anyone know?

PROT_EXEC ?
MAP_SHARED or MAP_PRIVATE. Choose one.

Got no file but still want to mmap? MAP_ANONYMOUS!

>mmap and fork

```
   fd = open("alice.txt", O_RDWR);

   char *ptr = mmap(NULL, 4096,
           PROT_READ|PROT_WRITE,
            MAP_FILE|MAP_SHARED, fd, 0);

if( (fork() ) ==0)
  strcpy(ptr, "The child wrote to the memory");
else {
  sleep(1); puts( ptr );
}
```

> What is RAID? Why is it necessary?

Making filesystems resilient:
RAID : "Redundant Array of Inexpensive Disks"

RAID Motivation

Mean Time to Failure (MTTF) ?
MTTF (disk array) = MTTF (single disk) / # disks
Adding more disks means that failures happen more frequently!

Simplest form: Mirroring "RAID 1"
All data is mirrored across two disks
Advantages:
Reads are faster, since both disks can be read in parallel
Higher reliability (of course)
Disadvantages:
Writes are slightly slower, since wait for both disks to do write
Doubles the cost of the storage system

RAID 3
Rather than mirroring, use parity codes
Given N bits {b1, b2, ..., bN}, the parity bit P is the bit {0,1} that yields an even number of "1" bits in the set {b1, b2, ..., bN, P}
Idea: If any bit in {b1, b2, ..., bN} is lost, can use the remaining bits (plus P) to recover it.
Where to store the parity codes? Add an extra "check disk" that stores parity bits

RAID 3 example
1. Read back data from other disks
2. Recalculate lost data from parity code
3. Rebuild data on lost disk

RAID 3 issues: performance

Terminology:
MTTF = mean time to failure
MTTR = mean time to repair
What is the MTTF of RAID?

Both RAID 1 and RAID 3 tolerate the failure of a single disk

RAID 5
Another approach: Interleaved check blocks ("RAID 5")
Rotate the assignment of data blocks and check blocks across disks
Avoids the bottleneck of a single disk for storing check data
Allows multiple reads/writes to occur in parallel (since different disks affected)

> A Planetary-sized Filesystem Case Study
Problem: Build a file system for Google

How do you make it resilient?
Reliable distributed storage
Issues
Failure is the common case
Google reports 2-10% of disks fail per year
Now multiply that by 60,000+ disks in a single warehouse...

Must survive failure of not just a disk, but failure of a rack of servers or even… a whole data center

How:
GFS 2001: Simple redundancy (2 or 3 copies of each file)

GFS 2010:
More efficient redundancy (analogous to RAID 3++)
Reed-Solomon codes with 1.5x redundancy
RS codes found in CDs, Space communication protocols