

# L P

→ canonical form

def primal LP  $\Pi = \max \langle c, x \rangle$   
s.t.  $Ax \leq b$   
 $x \geq 0$

$$c \in \mathbb{R}^n$$
$$A \in \mathbb{R}^{m \times n}$$

has dual  $\underline{\Pi}$

$\min \langle b, y \rangle$  ←  $m$  variables

s.t.  $A^T y \geq c$  ←  $n$  non-neg constraints  
 $y \geq 0$

$n$  variables.

$m$  non-neg constraints.

rmk dual is not in canonical form.

Thm → (weak duality)  
 $|\Pi| \leq |\underline{\Pi}|$  if both feasible.  
 $\max \langle c, x \rangle$        $\min \langle b, y \rangle$

Pf:  $x \geq 0$  feasible for  $\Pi \Rightarrow Ax \leq b$   
 $y \geq 0$  feasible for  $\underline{\Pi} \Rightarrow A^T y \geq c \Rightarrow y^T A \geq c^T \Rightarrow y^T A x \geq c^T x = \langle c, x \rangle$   
 $\parallel$   $y^T b = \langle b, x \rangle$

Cor dual is unbounded  $\Leftrightarrow (|\underline{\Pi}| = -\infty \Rightarrow \text{primal infeasible})$

primal is unbounded  $\Leftrightarrow (|\Pi| = \infty \Rightarrow \text{dual infeasible})$

### Thm (strong duality)

If  $\Pi$  feasible and bounded, then  $\mathcal{L}$  feasible and bounded and  $|\Pi| = |\mathcal{L}|$

### Thm (weak duality)

$$|\Pi| = \max_{\substack{Ax = b \\ Ax' \leq b' \\ x \geq 0}} \langle c, x \rangle \leq \min_{\substack{A^T y + (A')^T z \geq c \\ z \geq 0 \\ y \text{ unrestricted}}} \langle b, y \rangle + \langle b', z \rangle$$

Pf:

$$\begin{aligned} \langle c, x \rangle &= x^T c \leq x^T (A^T y + (A')^T z) \\ &= \underbrace{x^T A^T y}_{b^T} + \underbrace{x^T (A')^T z}_{\leq (b')^T} \end{aligned}$$

TODO

Now, back to maxflow, mincut

prop:  $\text{maxflow} = \max \sum_{e: s \rightarrow \circ} f_e - \sum_{e: \circ \rightarrow s} f_e$   
 s.t.  $f_{\text{out}} - f_{\text{in}} = 0$ ,  $f_e \leq c_e$ ,  $f_e \geq 0$

has dual  $\min \sum_e c_e x_e$   
 s.t.  $\begin{cases} d_v \leq d_u + x_e, & e: u \rightarrow v, \forall e. \\ d_s = 0, d_t = 1, & x_e \geq 0 \forall e. \end{cases}$

Pf: Rewrite maxflow LP

$$\begin{array}{l} \max \quad -F_t \\ d_s \left( \sum_{e: s \rightarrow \circ} f_e - \sum_{e: \circ \rightarrow s} f_e - F_s = 0 \right) \\ d_t \left( \sum_{e: t \rightarrow \circ} f_e - \sum_{e: \circ \rightarrow t} f_e - F_t = 0 \right) \\ d_v \left( \sum_{e: v \rightarrow \circ} f_e - \sum_{e: \circ \rightarrow v} f_e = 0 \right) \\ x_e \left( \begin{array}{l} f_e \leq c_e \\ f_e \geq 0 \end{array} \right) \end{array} \left. \vphantom{\begin{array}{l} \max \\ d_s \\ d_t \\ d_v \\ x_e \end{array}} \right\} \begin{array}{l} \text{Sum} \\ \text{up} \\ \Rightarrow -F_t \leq \min d_s \cdot 0 + d_t \cdot 0 + \sum_v d_v \cdot 0 + \sum_e c_e x_e \\ \text{s.t. } d_s = 0, d_t = 1, \underbrace{d_u - d_v + x_e}_{\geq 0} \geq 0, e: u \rightarrow v \\ \parallel \\ d_u \geq d_v + x_e \end{array}$$

Cor Strong duality  $\Rightarrow$  max flow = mincut

rmk When use FF, we need integral capacity to prove

Now, we prove the equality works for any capacity.