

# FLOW

flow through network?

def Directed graph  $G = (V, E)$  is simple,

{ no self loop  
no isolated vertex  
no parallel edges.

Will allow antiparallel.

def Capacitated graph.

simple directed graph

w/ edge capacity  $(C_e)_{e \in E}$ ,  $C_e \in \mathbb{N}$

Notation

for  $s, t \in V$ ,  $(s, t)$  flow in  $G$  is  $f = (f_e)_{e \in E}$ ,  $f_e \in \mathbb{R}_{\geq 0}$

$0 \leq f_e \leq C_e, \forall e.$

Def Conservation

$$f^{in}(v) = \sum_{e: u \rightarrow v} f_e$$

$$f^{out}(v) = \sum_{e: v \rightarrow u} f_e$$

$$f(v) = f^{out}(v) - f^{in}(v)$$

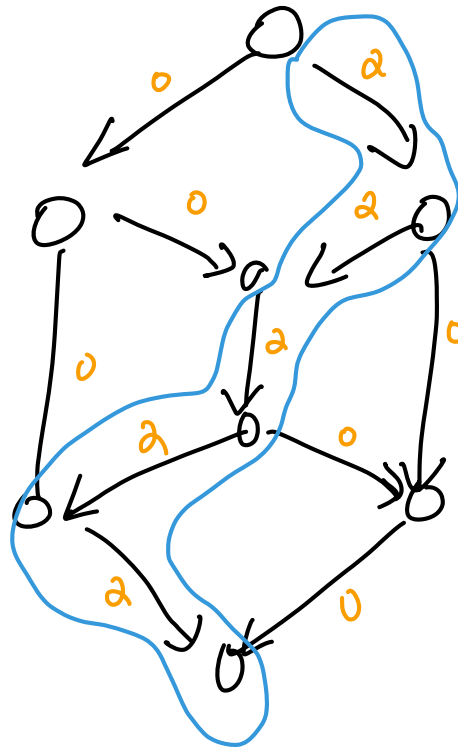
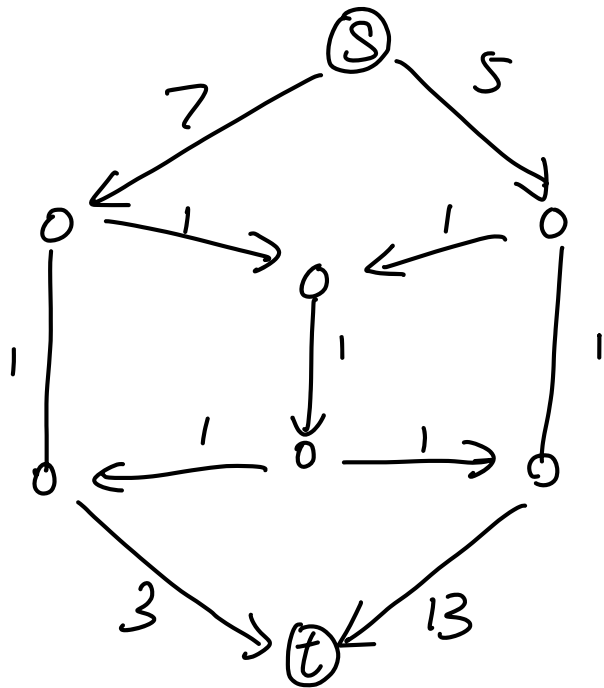
net flow:  $\forall v \in V \setminus \{s, t\}, f(v) = 0$

value of flow:  $|f| = f(s)$

Max flow is compute  $\text{arg max}_{f(s,t) \text{ flow}} |f|$

# Example

Capacity



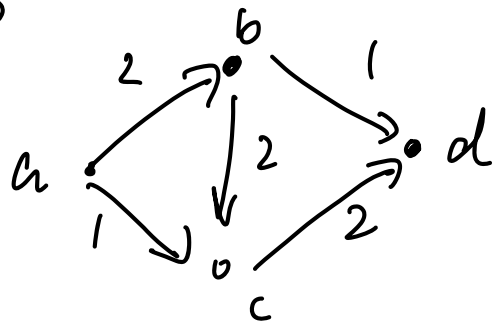
If  $\alpha = \pi - 3$ ,  
value of flow is  $\alpha = \pi - 3$

Cor max flow can be computed

Pf sketch: only finite # of integer flow.

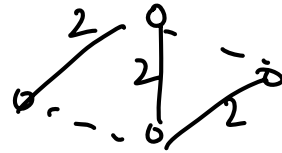
Q efficiency?   
 / DP ?   
 \ greedy ?

Greedy ?



idea: push along s-t path

Local max: one path  $a \rightarrow b \rightarrow d$  has capacity 2. So  $|f| = 2$



Global max:



$|f| = 3$ .

We can also push backwards.

def

$G^f = (V^f, E^f)$  residue graph

-  $V^f = V$

-  $E^f = \{e: e \in E, f_e < c_e\}$  ← forward edge

$\cup \{ \overline{-e}: e \in E, 0 < f_e \}$  ← backward edge

↑  
reversed

Then residue capacity,  $e \in E^f$ ,  $(C^f)_e = \begin{cases} c_e - f_e > 0 & e \text{ forward} \\ f_{-e} > 0 & e \text{ backward.} \end{cases}$

def